

Steiner triple systems $S(2^m - 1, 3, 2)$ of 2-rank $r \leq 2^m - m + 1$: construction and properties

D.V. Zinoviev, V.A. Zinoviev

A.A. Kharkevich Institute for Problems of Information Transmission, Moscow,
Russia

ACCT2012 Pomorie, Bulgaria, June 15-21, 2012

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3 New Construction

4 New Construction

A Steiner system $S(v, k, t)$ is a pair (X, B) , X is a v -set (i.e. $|X| = v$) and B – the collection of k -subsets of X (called blocks) such that every t -subset (of t elements) of X is contained in exactly one block of B .

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Present a Steiner system $S(v, 3, 2)$ ($S(v, 4, 3)$) by the binary incidence matrix (rather a set of rows). It is a binary constant weight code C of length v , blocks of B are codewords.

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Now, we enumerate $S(v, 3, 2)$, where $v = 2^m - 1$, of rank $2^m - m + 1$ ($\min + 2$).

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Let $J(v) = \{1, \dots, v\}$ be the coordinate set of S_v . Set $u = (v - 3)/4$. Define the subsets J_i of $J(v)$:

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$$J(v) = J_1 \cup \dots \cup J_u \cup J_{u+1}.$$

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For a given code $(3, 2, 16)_4$ -code L , define the constant weight
 $(12, 3, 4, 16)$ -code $C(L)$:

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For $x \in E^u$ with $\text{supp}(\mathbf{x}) = \{j_1, j_2, j_3\}$, we define a
 $(4u, 3, 4, 16)$ -code

$$C(L; \mathbf{x}) = C(L; j_1, j_2, j_3) = \{(\mathbf{c}_1, \dots, \mathbf{c}_u) : (\mathbf{c}_{j_1}, \mathbf{c}_{j_2}, \mathbf{c}_{j_3}) \in C(L)\},$$

and $\mathbf{c}_i = (0000)$ if $i \neq j_1, j_2, j_3$ (i.e. insert 3 blocks into u blocks).

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- $S^{(2,1)} = S_{v-2}^{(2,1)} \cup S_{v-1}^{(2,1)} \cup S_v^{(2,1)}$ is the set of words $\{\mathbf{c}\}$, $\text{supp}(\mathbf{c}) = \{j_1, j_2, j_3\}$, $j_1, j_2 \in J_i$, and $j_3 \in J_{u+1}$. The set $S_{v-2}^{(2,1)}$:

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Define (split 6 words of weight 2 into 3 pairs):

$$V(1) = \{(1100), (0011)\}, V(2) = \{(1010), (0101)\},$$

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■ $S^{(3)} = \{\mathbf{c} = (0 \dots 0111) : \text{supp}(\mathbf{c}) = J_{u+1}\}.$

Theorem 1.

Let $S_u = S(u, 3, 2)$ be a Steiner system and $\mathbf{c}^{(s)}$, $s = 1, 2, \dots, k$ its words, $k = u(u-1)/6$. Let $S^{(1,1,1)}$, $S^{(2,1)}$ and $S^{(3)}$ be the sets, obtained by our construction, based on the families of $(3, 2, 16)_4$ -codes L_1, L_2, \dots, L_k and the constant weight $(4, 2, 4, 2)$ -codes $V(1)$, $V(2)$ and $V(3)$. Set

$$S = S^{(1,1,1)} \cup S^{(2,1)} \cup S^{(3)}.$$

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$$S = S^{(1,1,1)} \cup S^{(2,1)} \cup S^{(3)}.$$

Then, for any choice of the codes L_1, L_2, \dots, L_k , the set S is the Steiner triple system $S_v = S(v, 3, 2)$ of order $v = 4u + 3$ with rank

$$v - (u - \text{rk}(S_u)) - 2 \leq \text{rk}(S_v) \leq v - (u - \text{rk}(S_u)).$$

A system $S_u = S(u, 3, 2)$ of order $u = 2^l - 1$ is called boolean if its rank is $u - l$, i.e. it is formed by the codewords of weight 3 of the linear Hamming code of length u .

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Suppose $S_v = S(v, 3, 2)$ is a Steiner system of order $v = 2^m - 1 = 4u + 3$. Suppose that its rank not greater than $v - m + 2$.

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Then this system S_v is obtained from the boolean Steiner triple system $S_u = S(u, 3, 2)$ of order $u = 2^{m-2} - 1$ using our construction, described above.

Theorem 3.

The following is true:

- *Let $m \geq 4$ and $v = 2^m - 1 \geq 15$. Set $u = (v - 3)/4$ and $k = u(u - 1)/6$. Then, the number M_v of different Steiner triple systems $S(v, 3, 2)$ of order v , whose rank is not greater than $v - m + 2$, and the fixed dual code \mathcal{A}_m , is equal to*

$$M_v = (2^6 \cdot 3^2)^k \times (6)^u, \quad k = u(u - 1)/6.$$

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$$M_v = (2^6 \cdot 3^2)^k \times (6)^u, \quad k = u(u - 1)/6.$$

- The overall number $M_v^{(o)}$ of different Steiner triple systems $S(v, 3, 2)$, whose rank $\leq v - m + 2$, is equal to

$$M_v^{(o)} = \frac{v! \cdot (2^6 \cdot 3^2)^k \cdot (6)^u}{(u(u - 1)(u - 2) \cdots (u + 1)/2) \cdot (4!)^u \cdot 3!}.$$

A system $S(v, 3, 2)$ of order $v = 2^m - 1$ is called *Hamming*, if it can be embedded into a binary non-linear perfect $(v, 3, 2^{v-m})$ -code (denoted by H_v), i.e. if it is the set of words of weight 3 of the code H_v , which contains the zero codeword.

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Theorem 4.

Any Steiner triple system $S_v = S(v, 3, 2)$ of order $v = 2^m - 1$ and rank $\text{rk}(S_v) \leq 2^m - m + 1$ is a Hamming system.