

Two types of upper bounds on the smallest size of a complete arc in the plane $PG(2, q)$

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Outline

- 1 Introduction
- 2 Bound with a decreasing degree of $\ln q$ (estimates of computer results)
- 3 Probabilistic bounds (uniform-distribution-assumption+theoretical way)

INTRODUCTION NOTATION

$PG(2, q) \Rightarrow$ projective **space** of dimension 2 over Galois field F_q

n -**arc** \Rightarrow a set of n points no three of which are collinear
a line meeting an arc \Rightarrow **tangent** or **bisecant**

bisecant \Rightarrow a line intersecting an arc in **two** points

a **point** A of $PG(2, q)$ is **covered** by an arc \Rightarrow
 A lies on a **bisecant** of the arc

complete arc \Rightarrow **all points** of $PG(2, q)$ are covered
by bisecants of the arc
 \Rightarrow one may not add a new point to a complete arc

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$t_2(2, q) \Rightarrow$ the smallest size of a complete arc in $PG(2, q)$

HARD OPEN CLASSICAL PROBLEM: 1950 \rightarrow
exact value or upper bound on $t_2(2, q)$

$\bar{t}_2(2, q) \Rightarrow$ the smallest known size of a complete arc in $PG(2, q)$
including computer search

$$t_2(2, q) \leq \bar{t}_2(2, q)$$

exact values of $t_2(2, q)$ for $q \leq 32$

for $q = 31, 32$ D. Bartoli, A. Milani, S. Marcugini, F. Pambianco,
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KNOWN BOUNDS on $t_2(2, q)$

lower bound

for $q = p^3$ O. Polverino **1999**

$$t_2(2, q) > \begin{cases} \sqrt{2q} + 1 & \forall q \\ \sqrt{3q} + \frac{1}{2} & q = p^h, h = 1, 2, 3 \end{cases}$$

CDL-bound \Rightarrow upper bound with a constant degree of $\log q$ DDL-bound \Rightarrow upper bound with a decreasing degree of $\log q$

DDL-bounds are more convenient

KNOWN BOUNDS with a DEGREE of LOGARITHM

$$\text{CDL-bound } t_2(2, q) \leq d\sqrt{q} \log^c q, \quad c \leq 300$$

c, d constants independent of q J.H. Kim, V. Vu 2003

theoretical bound; probabilistic methods

$$\text{CDL-bound } t_2(2, q) < \sqrt{q} \ln^{0.73} q$$

$$\text{DDL-bound } t_2(2, q) < 0.7\sqrt{q}(\ln q)^{\frac{1}{\ln q} + 0.78}$$

estimates of computer results for all $q \leq 13627$

D. Bartoli, A.A. Davydov, G. Faina, S. Marcugini, F. Pambianco 2005–2013

small arcs + BDFMP = big love

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NEW DDL-bound with a DECREASING DEGREE of $\ln q$

Theorem

$$\text{DDL-bound} \quad t_2(2, q) < 0.6\sqrt{q}(\ln q)^{\varphi(q)}$$

$$\varphi(q) = \frac{1.5}{\ln q} + 0.8$$

all $q \leq 49727$, all prime $q \leq 89113$

100 sporadic prime q with $90001 \leq q \leq 350003$

estimates of computer results in HUGE region
ALL prime q WITHOUT GAPS

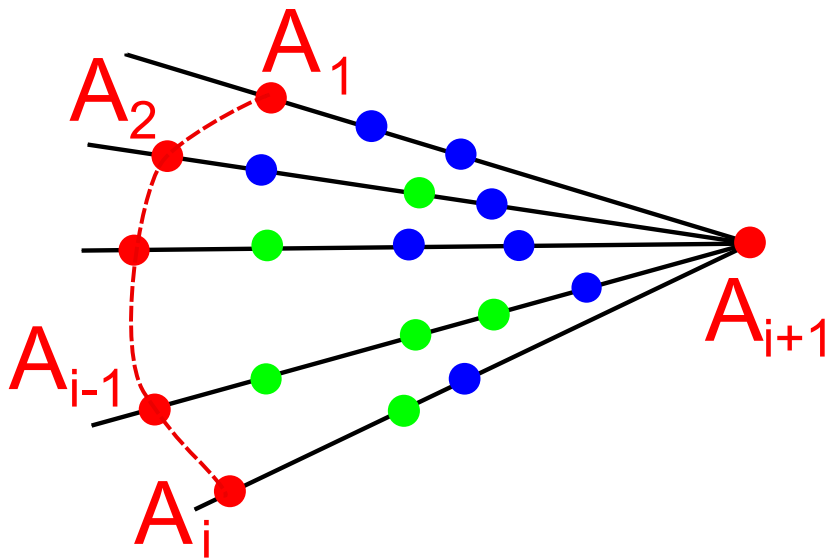
Randomized greedy algorithms

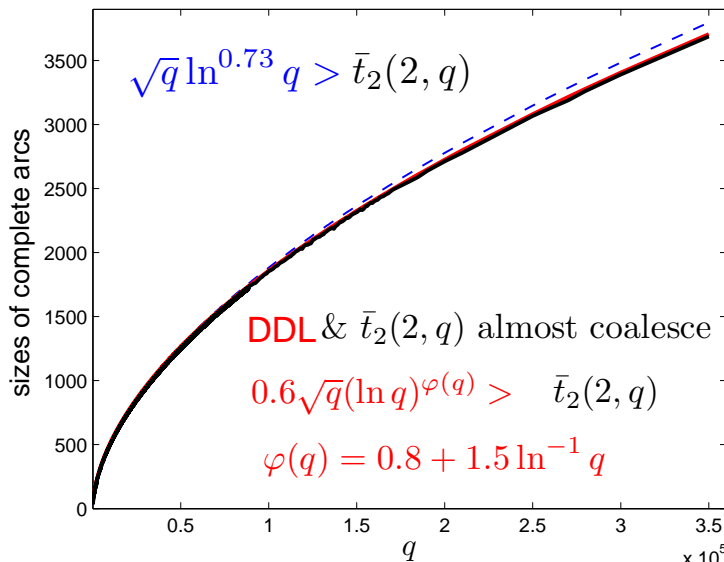
New arcs have been obtained using resources of Multipurpose Computing Complex of National Research Centre “Kurchatov Institute”

A greedy algorithm is an algorithm that makes the *locally optimal choice* at each stage with the hope of finding a global optimum or, at least, a global “good” solution.

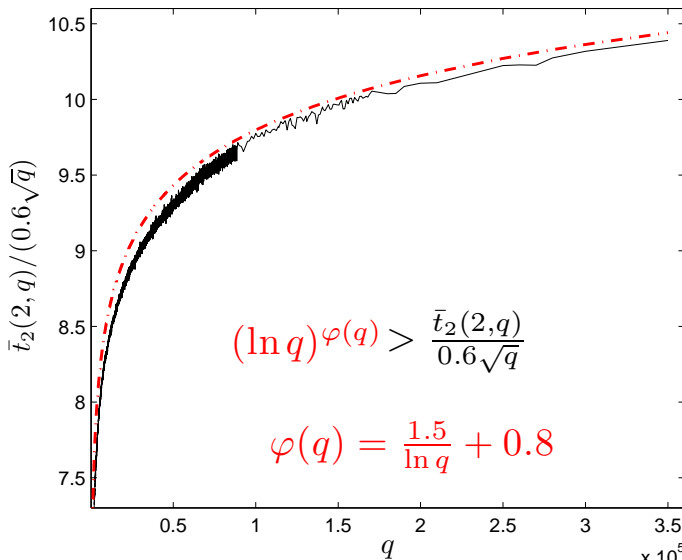
A randomized greedy algorithm executes some stages in a random manner without the local optimum.

One step of greedy algorithm



DECREASING-DEGREE-LOGARITHM-bound $>$ Arc's Size

$$(\ln q)^{\varphi(q)} > \bar{t}_2(2, q)/(0.6\sqrt{q})$$



Assumption of a uniform distribution

$\theta = q^2 + q + 1 \Rightarrow$ the number of points in $PG(2, q)$

$NONcov_i \Rightarrow$ the number of noncovered points of $PG(2, q)$ after the i -th step of the algorithm.

Assumption 1

*Covered and noncovered points of the plane are distributed uniformly. The **proportion** $\frac{NONcov_i}{\theta}$ of noncovered points **is** the **probability** that a random point of the plane is noncovered on the $i + 1$ -th step.*

PROBABILISTIC UPPER BOUNDS

Theorem

Under *Assumption 1*,

$$2\sqrt{q} \ln^{0.5} q > B(q) > t_2(2, q).$$

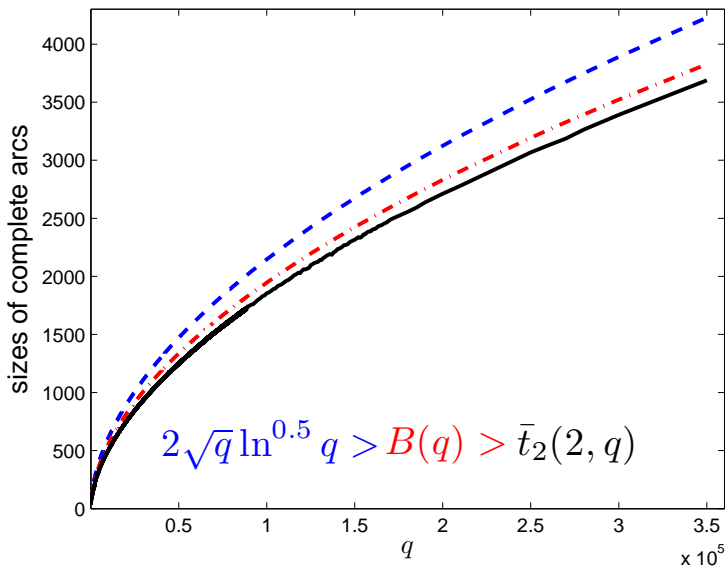
$B(q) = x$ is the solution of equation

$$\prod_{i=1}^x \left(1 - \frac{i}{q}\right) = \frac{1}{\theta} \left(1 + \sum_{j=1}^x \prod_{i=j}^x \left(1 - \frac{i}{q}\right)\right)$$

assumption on uniform distribution + theoretical way

$2\sqrt{q} \ln^{0.5} q \Rightarrow$ rough estimate of $B(q)$

Probabilistic upper bounds



everything will be fine

All above mentioned new upper bounds are confirmed in huge region of q .

Conjecture

All above mentioned new upper bounds hold for all q .

Thank you Spasibo
Mille grazie
Premnogo blagodarya
!'Muchas gracias
Toda raba
Merci beaucoup
Dankeschön
Dank u wel
Domo arigato