

On the Extendability of Quasidivisible Optimal Arcs ¹

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Abstract. We prove a new sufficient condition for the extendability of Griesmer arcs with certain parameters.

1 Preliminaries

It is a well-known fact that adding a parity check bit to all words of a binary $[n, k, d]$ -code with d odd increases the minimum distance, i.e. the resulting code has parameters $[n + 1, k, d + 1]$. This result has been generalized by Hill and Lizak in [4, 5]. They showed that if all weights in an $[n, k, d]_q$ code are congruent to 0 or $d \pmod{q}$, with $(d, q) = 1$, then it can be extended to an $[n + 1, k, d + 1]_q$ code. It turned out that this fact has a natural explanation in terms of blocking sets with respect to hyperplanes. It was proved in [6, 9] that the result of Hill and Lizak can be obtained from the well-known Bose-Burton theorem for blocking sets in $\text{PG}(k - 1, q)$. A further generalization along these lines was done in [7] by exploiting a result of Beutelspacher and Heim on the size of the minimal non-trivial (i.e. not containing a hyperplane) blocking set in a finite projective geometry.

In a series of papers, Maruta obtained various results [9–11] on the extendability of linear codes. He introduced the notion diversity of a linear code with spectrum (A_i) as the pair (Φ_0, Φ_1) , where

$$\Phi_0 = \frac{1}{q-1} \sum_{q|i, i \neq 0} A_i, \quad \Phi_1 = \frac{1}{q-1} \sum_{i \neq 0, d(q)} A_i,$$

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and proved that for various values of the diversity the code is indeed extendable. In particular, he showed that a linear $[n, k, d]$ code over \mathbb{F}_q , with $q \geq 5$ odd, having all non-zero weights congruent to $-2, -1$, and 0 modulo q is extendable.

In this note, we introduce the notion of t -quasidivisibility modulo q . We define a special arc $\tilde{\mathcal{K}}$ in the dual geometry $\widehat{\text{PG}}(k-1, q)$ and relate the extendability property for \mathcal{K} with the existence of a hyperplane in the support of $\tilde{\mathcal{K}}$. The main theorem in this note states that every t -quasidivisible Griesmer arc with divisor q , $t < \sqrt{q}$, and an additional numerical condition on the parameters, is t -extendable.

2 Basic definitions

Let \mathcal{P} be the set of points of the projective geometry $\text{PG}(k-1, q)$. Every mapping $\mathcal{K} : \mathcal{P} \rightarrow \mathbb{N}_0$ from the pointset of the geometry to the non-negative integers is called a multiset in $\text{PG}(k-1, q)$. This mapping is extended additively to the subsets of \mathcal{P} : for every $\mathcal{Q} \subseteq \mathcal{P}$, $\mathcal{K}(\mathcal{Q}) = \sum_{P \in \mathcal{Q}} \mathcal{K}(P)$. The integer $n := \mathcal{K}(\mathcal{P})$ is called the cardinality of \mathcal{K} . For every set of points $\mathcal{Q} \subset \mathcal{P}$ we define its characteristic (multi)set $\chi_{\mathcal{Q}}$ by

$$\chi_{\mathcal{Q}}(P) = \begin{cases} 1 & \text{if } P \in \mathcal{Q}, \\ 0 & \text{otherwise.} \end{cases}$$

Multisets can be viewed as arcs or as blocking sets. A multiset \mathcal{K} in $\text{PG}(k-1, q)$ is called an (n, w) -multiarc (or simply (n, w) -arc) if (1) $\mathcal{K}(\mathcal{P}) = n$, (2) $\mathcal{K}(H) \leq w$ for every hyperplane H , and (3) there exists a hyperplane H_0 with $\mathcal{K}(H_0) = w$. Similarly, a multiset \mathcal{K} in $\text{PG}(k-1, q)$ is called an (n, w) -blocking set with respect to the hyperplanes (or (n, w) -minihyper) if (1) $\mathcal{K}(\mathcal{P}) = n$, (2) $\mathcal{K}(H) \geq w$ for every hyperplane H , and (3) there exists a hyperplane H_0 with $\mathcal{K}(H_0) = w$.

An (n, w) -arc \mathcal{K} in $\text{PG}(k-1, q)$ is called t -extendable, if there exists an $(n+t, w)$ -arc \mathcal{K}' in $\text{PG}(k-1, q)$ with $\mathcal{K}'(P) \geq \mathcal{K}(P)$ for every point $P \in \mathcal{P}$. Arcs that are 1-extendable are called simply extendable arcs. An arc is said to be complete if it is not extendable. Similarly, an (n, w) -blocking set \mathcal{K} in $\text{PG}(k-1, q)$ is called reducible, if there exists an $(n-1, w)$ -blocking set \mathcal{K}' in $\text{PG}(k-1, q)$ with $\mathcal{K}'(P) \leq \mathcal{K}(P)$ for every point $P \in \mathcal{P}$. A blocking set is called irreducible if it is not reducible.

Given a multiset \mathcal{K} in $\text{PG}(k-1, q)$, we denote by a_i the number of hyperplanes H with $\mathcal{K}(H) = i$. The sequence (a_i) is called the spectrum of \mathcal{K} . A multiset \mathcal{K} with $\mathcal{K}(\mathcal{P}) = n$ and spectrum (a_i) is said to be divisible with divisor $\Delta > 1$ if $a_i = 0$ for all $i \not\equiv n \pmod{\Delta}$. \mathcal{K} is called t -quasidivisible with divisor $\Delta > 1$ (or t -quasidivisible modulo Δ) if $a_i = 0$ for all $i \not\equiv n, n+1, \dots, n+t \pmod{\Delta}$. An easy corollary of Hill and Lizak's result is that every 1-quasidivisible arc with divisor q is extendable. Maruta's

theorem from [11] claims that for q odd every 2-quasidivisible arc with divisor q is extendable.

3 Extendability of Griesmer Arcs

There exists an one-to-one correspondence between the classes of isomorphic $[n, k, d]_q$ codes and the classes of projectively equivalent $(n, n-d)$ -arcs in $\text{PG}(k-1, q)$ [3]. With every multarc \mathcal{K} we can associate many isomorphic linear codes. Fix one of them and denote it by $C_{\mathcal{K}}$. If $C_{\mathcal{K}}$ is a Griesmer code then we call \mathcal{K} a Griesmer arc.

Let \mathcal{K} be a t -quasidivisible (n, w) -arc with divisor q in $\Sigma = \text{PG}(k-1, q)$, $t < q$. Set $d = n - w$. This is a typical situation when one investigates the existence problem for Griesmer arcs with given parameters.

Define a new multiset $\tilde{\mathcal{K}}$ in the dual geometry $\tilde{\Sigma}$ by

$$\tilde{\mathcal{K}} : \begin{cases} \mathcal{H} & \rightarrow \mathbb{N}_0 \\ H & \rightarrow \tilde{\mathcal{K}}(H) = n + t - \mathcal{K}(H) \pmod{q}. \end{cases} \quad (1)$$

In other words, hyperplanes of multiplicity congruent to $n + a \pmod{q}$ become $(t - a)$ -points in the dual geometry. The following result is straightforward.

Theorem 1. *Let \mathcal{K} be an (n, w) -arc in $\Sigma = \text{PG}(k-1, q)$ that is t -quasidivisible modulo q with $t < q$. Let $\tilde{\mathcal{K}}$ be defined by (1). If*

$$\tilde{\mathcal{K}} = \sum_{i=1}^c \chi_{\tilde{H}_i} + \tilde{\mathcal{K}}'$$

for some multiset $\tilde{\mathcal{K}}'$ and c not necessarily different hyperplanes $\tilde{H}_1, \dots, \tilde{H}_c$ in the dual geometry $\tilde{\Sigma}$, then \mathcal{K} is c -extendable. In particular, if $\tilde{\mathcal{K}}$ contains a hyperplane in its support then \mathcal{K} is extendable.

By Theorem 1, the extendability of t -quasidivisible arcs is linked with the structure of the multiset $\tilde{\mathcal{K}}$ defined in the dual geometry. It turns out that this multiset is highly divisible.

Theorem 2. *Let \tilde{S} be subspace of $\tilde{\Sigma}$ of dimension at least 1. Then*

$$\tilde{\mathcal{K}}(\tilde{S}) \equiv t \pmod{q}.$$

By this theorem, the multiset $\tilde{\mathcal{K}}$ has the following properties:

- the multiplicity of each point is at most t ;
- the multiplicity of each subspace of dimension r , where $1 \leq r \leq k-1$, is at least tv_r .

Here we use the conventional notation $v_r = (q^r - 1)/(q - 1)$. Let us note that in the general case we do not know the cardinality of $\tilde{\mathcal{K}}$.

4 A Theorem on the Extendability of Griesmer Arcs

Consider a Griesmer t -quasidivisible arc \mathcal{K} , $t < q$, with parameters (n, w) in $\text{PG}(k - 1, q)$. Set $d = n - w$ and let $C_{\mathcal{K}}$ be a linear code associated with \mathcal{K} . The code $C_{\mathcal{K}}$ has parameters $[n, k, d]_q$, where d can be written as

$$d = sq^{k-1} - \sum_{i=0}^{k-2} \varepsilon_i q^i, \quad 0 \leq \varepsilon_i < q. \tag{2}$$

Then we have $\lceil d/q^j \rceil = sq^{k-j-1} - \sum_{i=j}^{k-2} \varepsilon_i q^i$, which implies

$$n = sv_k - \sum_{i=0}^{k-2} \varepsilon_i v_{i+1}. \tag{3}$$

Let us note that with this notation $t = \varepsilon_0$ since $n + \varepsilon_0 \equiv w_1 \pmod{q}$. Denote by w_j the maximal multiplicity of a subspace S of codimension j of $\text{PG}(k - 1, q)$: $w_j = \max_{\text{codim } S=j} \mathcal{K}(S)$, $j = 1, \dots, k - 1$. We have

$$w_j = \sum_{i=j}^{k-1} \lceil \frac{d}{q^i} \rceil = sv_{k-j} - \sum_{i=j}^{k-2} \varepsilon_i v_{i-j}. \tag{4}$$

By convention $w_0 = n$.

In the next lemmas we establish some important properties of the arc $\tilde{\mathcal{K}}$. The proofs of the first two lemmas use simple counting arguments.

Lemma 1. *Let \mathcal{K} be a t -quasidivisible $(n, n - d)$ -Griesmer arc with d given by (2). Let S be a subspace of codimension 2 contained in the hyperplane H_0 with $\mathcal{K}(H_0) = w_1 - aq$ where $a \geq 0$ is an integer.*

- (i) *If $\mathcal{K}(S) = w_2 - a - b$, $0 \leq b \leq t - 2$, then $\tilde{\mathcal{K}}(\tilde{S}) \leq t + bq$;*
- (ii) *If $\mathcal{K}(S) = w_2 - a - b$, $b \geq t - 1$, then $\tilde{\mathcal{K}}(\tilde{S}) \leq t + (t - 1)q$.*

Lemma 2. *Let \mathcal{K} and $\tilde{\mathcal{K}}$ be as in Lemma 1. Let T be a subspace of codimension 3 in $\text{PG}(k - 1, q)$ with $\mathcal{K}(T) = w_3$. Then*

$$\tilde{\mathcal{K}}(\tilde{T}) \leq t(q + 1) + \varepsilon_1 q.$$

The next lemma shows that for certain parameters the arc $\tilde{\mathcal{K}}(\tilde{T})$ is minimal, i.e. a sum of lines.

Lemma 3. *Let \mathcal{K} be t -quasidivisible Griesmer (n, w) -arc in $\text{PG}(k - 1, q)$, $q \geq 3$ with $d = n - w$ given by (2). Let $\tilde{\mathcal{K}}$ be defined by (1). Let further $\varepsilon_0, \varepsilon_1 \leq \sqrt{q}$. For every maximal subspace T of codimension 3 in $\text{PG}(k - 1, q)$, i.e. a subspace with $\mathcal{K}(T) = w_3$, it holds*

$$\tilde{\mathcal{K}}(\tilde{T}) = t(q + 1).$$

Proof. We have that \tilde{T} is a plane in $\widetilde{\text{PG}}(k-1, q)$. By Lemma 2, $\tilde{\mathcal{K}}(\tilde{T}) \leq \varepsilon_0(q+1) + \varepsilon_1q$. Set $\tilde{\mathcal{K}}(\tilde{T}) = \varepsilon_0(q+1) + \varepsilon'_1q$, where $0 \leq \varepsilon'_1 \leq \varepsilon_1$.

Assume $\varepsilon'_1 > 0$. Set $\tilde{\mathcal{F}} = \tilde{\mathcal{K}}|_{\tilde{T}}$, i.e. $\tilde{\mathcal{F}}$ is the restriction of $\tilde{\mathcal{K}}$ to the plane \tilde{T} in the dual geometry. Define a dual plane arc \mathcal{F} to $\tilde{\mathcal{F}}$ by

$$\mathcal{F}(\tilde{L}) = i \quad \text{iff} \quad \tilde{\mathcal{F}}(L) = t + iq.$$

Denote by (A_i) the spectrum of $\tilde{\mathcal{F}}$. We have

$$\begin{aligned} \sum A_{t+iq} &= q^2 + q + 1 \\ \sum (t+iq)A_{t+iq} &= \varepsilon(q+1) + \varepsilon'_1q \end{aligned}$$

for some $\varepsilon'_1 \leq \varepsilon_1$. This implies $\sum_i A_{t+iq} = \varepsilon'_1(q+1) + \varepsilon_0$.

Now let us denote by B_i the number of lines L with $\tilde{\mathcal{F}}(L) = t + iq$ through a fixed point P of multiplicity $c \geq 0$. Then

$$\begin{aligned} \sum B_{t+iq} &= q + 1 \\ \sum (t+iq)B_{t+iq} &= (q+1)\varepsilon_0 + \varepsilon'_1q + cq \end{aligned}$$

which implies $\sum_i iB_{t+iq} = \varepsilon'_1 + c$. Hence \mathcal{F} is a $(\varepsilon'_1(q+1) + \varepsilon_0, \varepsilon'_1)$ -blocking set.

From $\varepsilon_0, \varepsilon_1 < \sqrt{q}$ and $q \geq 3$ we get that $\varepsilon_0 + \varepsilon'_1 < \sqrt{\varepsilon'_1q} + 1$ and, consequently, $\varepsilon'_1(q+1) + \varepsilon_0 < \varepsilon'_1q + \sqrt{\varepsilon'_1q} + 1$. By a well known result Ball [1] and De Beule-Storme-Metsch [2], \mathcal{F} contains a line. Going back to $\tilde{\mathcal{F}}$ this implies that all lines L_i in \tilde{T} through \tilde{P} have multiplicity at least $t + q = \varepsilon_0 + q$. Now we have

$$\begin{aligned} \varepsilon_0(q+1) + \varepsilon_1q \geq \tilde{\mathcal{K}}(\tilde{T}) &= \sum_{i=0}^q \tilde{\mathcal{K}}(L_i) - q\tilde{\mathcal{K}}(\tilde{P}) \\ &\geq (q+1)(\varepsilon_0 + q) - q\tilde{\mathcal{K}}(\tilde{P}) \\ &\geq \varepsilon_0(q+1) + q(q+1) - \varepsilon_0q. \end{aligned}$$

This implies $q+1 \leq \varepsilon_0 + \varepsilon_1 < 2\sqrt{q}$, i.e. $(\sqrt{q}-1)^2 < 0$, which is a contradiction. Therefore $\varepsilon'_1 = 0$ which proves the lemma. \square

Using induction on the dimension, we can extend Lemma 3 to subspaces of arbitrary dimension.

Lemma 4. *Let \mathcal{K} be t -quasidivisible Griesmer (n, w) -arc in $\text{PG}(k-1, q)$, $q \geq 3$ with $d = n - w$ given by (2). Let $\tilde{\mathcal{K}}$ be defined by (1). Let U be a subspace in $\text{PG}(k-1, q)$ of maximal multiplicity with $\text{codim } U = r$, $1 \leq r \leq k$ (if $\text{codim } U = k$, $U = \emptyset$). If $\varepsilon_0, \varepsilon_1, \dots, \varepsilon_{r-2} < \sqrt{q}$, then*

$$\tilde{\mathcal{K}}(\tilde{U}) = \varepsilon_0 v_{r-1}.$$

Now our main theorem becomes almost straightforward.

Theorem 3. *Let \mathcal{K} be a t -quasidivisible Griesmer $(n, n - d)$ -arc with d given by (2). Let $\varepsilon_0 = t, \dots, \varepsilon_{k-2} < \sqrt{q}$. Then \mathcal{K} is t -extendable.*

Proof. By Lemma 4, $\tilde{\mathcal{K}}$ is a (tv_{k-1}, tv_{k-2}) -minihyper. By corollary 3.5 from [8], every (tv_{k-1}, tv_{k-2}) minihyper in $\text{PG}(k-1, q)$ with $t \leq q - \frac{q}{p}$ is the sum of hyperplanes. Now the result follows from Theorem 1 since $t < \sqrt{q} < q - \frac{q}{p}$. \square

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