# Partitions and constant-value codes

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**Abstract.** We study the relationship between partitions of some integer a in GF(p) in unequal parts of size at most (p-1)/2, and binary vectors with so-called value a. In particular we investigate a group of transformations acting on the family  $A = \{A, A, ..., A\}$ , where A stands for the set of all vectors of value i.

### **1** Preliminaries

Let p be some odd prime. We shall study the partitions of positive integers consisting of unequal parts the size of which is at most (p-1)/2. It will be obvious that we can represent such partitions by binary vectors  $c = (c_1, c_2, \cdots, c_{(p-1)/2})$  of length (p-1)/2. Here,  $c_i = 1$  if and only if the partition contains a part of size i. We interpret all vectors as row vectors. The number of ones in such a vector c is called the weight of the partition and is denoted by |c|. It stands for the number of parts in the partition. Let c be some partition. We define

$$a = \sum_{j=1}^{(p-1)/2} jc_j \mod p$$
 (1)

and call a the value of c or val(c), with  $a \in \{0, 1, \dots, p-1\}$ . For a fixed value a, we collect all vectors having this value in a set  $A_a$  consisting of  $|A_a|$  binary vectors of length (p-1)/2. So, this set contains all "conventional" partitions of the integers  $a, a + p, a_2p, \cdots$  into unequal parts. We shall call such a set a *constant-value code*. We also introduce integers  $n_e$  and  $n_o$ , being the number of vectors in  $A_a$  with an even number of ones and an odd number, respectively. (We suppress the a-dependency of these integers in our notation). The complement of a partition c is defined as the partition corresponding to the vector c = c + 1, where 1 is the all-one vector of length (p - 2)/2. Since the value of 1 is equal to

$$L := (p^2 - 1)/8 \mod p$$
 (2)

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all vectors of a set  $A_a$  have a complement of the same value L - a. Hence, we can write  $A_a^c = A_{L-a}$  and we call  $A_a^c$  the complement of  $A_a$ . We also need the "value of the first halve of 1", defined by

$$K = 1 + 2 + \dots + [(p-1)/4] = (p^2 \mp 2p - 3)/32 \mod p, \tag{3}$$

for  $p = \mp 1 \mod 4$ . Consequently we have

$$L - 4K = (1 \pm p)/4 \mod p$$
 (4)

Finally, we introduce the number  $k \in GF(p)$ , defined by

$$2k = L = (p^2 - 1)/8 \tag{5}$$

as equality in GF(p). In order to deal with the sets  $A_a, a \in \{0, 1, \dots, p-1\}$ , we also introduce

$$N(p) = \begin{cases} \frac{2^{(p-1)/2} + 1}{p}, & p = \pm 3 \mod p; \\ \frac{2^{(p-1)/2} - 1}{p}, & p = \pm 1 \mod p. \end{cases}$$
(6)

### 2 A group of transformations

Let  $I = \{1, 2, \dots, (p-1)/2\}$  and let m be some integer with  $1 \le m \le p-1$ . We introduce index sets

$$I_1 = \{i : i \in I, mi \mod p \in I\}, \quad I_2 := I \setminus I_1 \tag{7}$$

and a permutation matrix P with elements

$$p_{ij} = 1, \ j = mi \mod p, \ i \in I_1, or \ j = -mi \mod p, \ i \in I_2$$

$$\tag{8}$$

while  $p_{i,j} = 0$  otherwise.

**Theorem 1.** Let l be the order of  $m \mod p$ . Then the matrix P defined by (8) represents a permutation on I consisting of (p-1)/l cycles of length l/2, for l is even, and of (p-1)/2l cycles of length l, for l is odd.

*Proof.* Consider the mapping  $\mathcal{P}: GF(p) \to GF(p), \mathcal{P} = ma$ . This mapping gives rise to a permutation of the elements of I in the following way. First,  $\mathcal{P}$  permutes the nonzero elements of GF(p) according to (p-1)/l cycles of length l. Next, we change all elements a in these cycles which are not in I into a' := a - p, and then omit the minus sign of a'. If -1 is in the same cycle as 1, which is the case for l is even, this cycle of length l/2, while all elements now

are in I. The same holds for all other cycles. If -1 and 1 are in different cycles of length l, which is the case for l is odd, then both cycles become identical after changing the minus signs. So, when omitting repeated cycles, we end up with a permutation of the elements of I as described in the theorem. For the matrix P the same holds. More precisely, this matrix represents the mapping  $\mathcal{P}^{-1}$ , modified by the above procedure.

Next, we define a translation vector  $t = (t_1, t_2, \cdots, t_{(p-1)/2})$ , with  $t_j = 1$  for  $j = mi \mod p$ , for  $i \in I_1$ , and  $t_j = 0$  otherwise. Furthermore, we consider the transformation  $T_m := GF(p)^{(p-1)/2} \to GF(p)^{(p-1)/2}$  defined by

$$T_m(c) = cP + t \tag{9}$$

**Theorem 2.** For each m,  $1 \le m \le p-1$ ,  $T_m$  induces a permutation  $\tau_m$ on the set  $A = A_0, A_1, \dots, A_{p-1}$  such that  $\tau_m(A_a) = A_b$ , with  $b = m(S_m - a)$ and  $S_m = \sum_{i \in I_1} i$ .

*Proof.* We shall determine the value w' of the vector  $b = T_m(a)$ , with val(a) = w. The components  $i \in I$  contribute  $\sum_{i \in I_1} mi(1-a_i)$  to w' and those in  $I_2$  yield  $\sum_{i \in I_2} (p - mia_i)$ . Hence, both contributions together and taken mod p, give  $w' = \sum_{i \in I_1} mi - \sum_{i \in I} mia_i = mS - mw$ .

Special cases

$$\begin{split} m &= 2 & I_1 = 1, 2, \cdots, [(p-1)/4], \ I_2 = I \setminus I_1, \\ t &= (0, 1, 0, 1 \cdots), \ w' = 2(S_2 - w) = 2(K - w); \\ m &= (p-1)/2 & I_1 = 1, 3, 5, \dots, \ I_2 = 2, 4, 6, \dots, t = (1, 0, 1, 0 \dots), \\ w' &= (p-1)/2 \cdot (S_{(p-1)/2} - w) = (p-1)/2 \cdot (L - K - w); \\ m &= p-1 & I_1 = \emptyset, \ I_2 = I, \ t = 0, P = E. \end{split}$$

Let  $w_{i,n}$  be the value of the set  $\tau_m^n(A_i)$ . The integers  $w_{i,n}$  satisfy in GF(p) the recurrence relation

$$w_{i,n} = m(S_m - w_{i,n-1}), \ w_{i,0} = i, \tag{10}$$

which has as solution

$$w_{i,n} = \frac{m}{m+1} S_m (1 - (-m)^n) + i(-m)^n.$$
(11)

The permutations  $\tau_m, 1 \leq m \leq p-1$ , generate a permutation group  $G_A$  on A.

#### Theorem 3.

- (i)  $G_A$  can be generated by a permutation  $\tau_{-\alpha}$ , where  $\alpha$  is a generator of  $GF(p)^*$ .
- (ii)  $G_A$  has one orbit  $A_k$  of size 1, whereas all other  $A_i$ ,  $i \neq k$ , are in one orbit of size p-1.

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*Proof.* Since  $\alpha$  generates the multiplicative group of GF(p), we can write  $m = \alpha^e$  for any  $m \in 1, 2, ..., p-1$ . The permutation  $\tau_m$  generates a subgroup of  $G_A$ . Equality (10) implies that  $\frac{m}{m+1}S_m$  has the same value for all m. Since  $S_1 = L$ , it follows that  $\frac{m}{m+1}S_m = \frac{L}{2}$ . Next, from (10) and (11) we have that  $w_{i,n} = i$  is equivalent to

$$(L/2 - i)(1 - (-m)^n) = 0$$
(12)

The only *i*-value which satisfies this equation is i = L/2 = k. So,  $A_k$  is invariant with respect to all transformations of  $G_A$ . Furthermore, it will be clear from (11), that the length of the orbit to which  $A_i, i \neq k$ , belongs under the action of  $\tau_m$ , is equal to the order of  $-m \mod p$ . So, if we take  $m = -\alpha$ , the orbit has length p - 1.

**Example.** For p = 11 we have the following data: L = 4, k = 3, K = 3. The family A of constant-value codes consists of the sets:

$$\begin{split} &A_0 = (0,0,0,0,0), (0,1,0,1,1), (1,1,1,0,1) \quad A_1 = (1,0,0,0,0), (0,0,1,1,1), (1,1,0,1,1) \\ &A_2 = (0,1,0,0,0), (1,0,1,1,1) \\ &A_3 = (0,0,1,0,0), (1,1,0,0,0), (0,1,1,1,1) \quad A_4 = (0,0,0,1,0), (1,0,1,0,0), (1,1,1,1,1) \\ &A_5 = (0,0,0,0,1), (1,0,0,1,0), (0,1,1,0,0) \quad A_6 = (1,0,0,0,1), (0,1,0,1,0), (1,1,1,1,0,0) \\ &A_7 = (0,1,0,0,1), (0,0,1,1,0), (1,1,0,1,0) \quad A_8 = (0,0,1,0,1), (1,0,1,1,0), (1,1,0,0,1) \\ &A_9 = (0,0,0,1,1), (0,1,1,1,0), (1,0,1,0,1) \quad A_{10} = (1,0,0,1,1), (0,1,1,0,1), (1,1,1,1,0) \end{split}$$

In this case, 2 generates the multiplicative group of the relevant field, i.e.  $GF(11)^*$ . So, according to Theorem 3 the transformation  $\tau_{-2} = \tau_9$  is a generator of  $G_A$ , and it acts transitively on the family  $A_i | i \neq k$ . In order to apply Theorem 2, we obtain  $I_1 = 3, 4, 5$ , and hence  $S_9 = 3+4+5=1 \mod 11$ . Indeed, the relations  $\tau_9(A_a) = A_b$  and b = 9(1-a) provide us with the transformations:

 $A_2 \rightarrow A_2, \ A_0 \rightarrow A_9 \rightarrow A_5 \rightarrow A_8 \rightarrow A_3 \rightarrow A_4 \rightarrow A_6 \rightarrow A_{10} \rightarrow A_7 \rightarrow A_1 \rightarrow A_0$ 

## **3** Constructing $A_{i+1}$ from $A_i$

Next, we shall discuss a method to transform a vector  $a \in A_i$  into a vector  $b \in A_{i+1}$ . For the sake of convenience we assume that 2 is a generator of  $GF(p)^*$ . So, the matrix P in (8) corresponds to a (p-1)/2-cycle which we denote by

$$d := (d_1(=1), d_2, \dots, d_{(p-1)/2}), \quad d_i \in I$$
(13)

Corresponding to (13) we define a binary vector p of length (p-1)/2, such that its *i*-th component is equal to the parity of the number of  $d_j$ , j < i, which are in  $I_2$ .

Now, let a be a binary vector representing some partition, and let val(a) = i. We define a translation vector t as follows. If  $a_{d_j} \neq p_j$ ,  $1 \leq j < k$ , and  $a_{d_k} = p_k$  for some k,  $1 \le k \le (p-1)/2$ , we put  $t_{d_j} = 1$ , whereas all other components are zero. Formally, we can obtain t by

$$t = (1, \dots, 1, 0, \dots, 0)Q \tag{14}$$

where the vector at the rhs contains k ones followed by (p-1)/2 - k zeros, while the transformation matrix Q has elements  $q_{i,j} = 1$  if  $j = d_i$  and  $q_{i,j} = 0$ otherwise.

Theorem 4.

- (i) If  $a \in A_i$ , then  $b = a + t \in A_{i+1}$ , unless  $a = a_0 := p^c Q$ ;
- (ii) For  $p = \pm 3 \mod 8$ , the translation in (i) gives one-to-one mappings  $A_i \rightarrow A_{i+1}, \forall i \in GF(p) \setminus \{k-1,k\}, A_{k-1} \setminus \{a_0\} \rightarrow A_k \text{ and } A_k \rightarrow A_{k+1} \setminus \{a_0^c\};$
- (iii) For  $p = \pm 1 \mod 8$ , the translation in (i) gives one-to-one mappings  $A_i \to A_{i+1}, \forall i \in GF(p) \setminus \{k-1,k\}, A_{k-1} \to A_k \setminus \{a_0^c\} \text{ and } A_k \setminus \{a_0\} \to A_{k+1}$ .

*Proof.* We only have to take into account the change in the contribution to val(a) due to the components  $a_{d_1}, \ldots, a_{d_k}$ . These contribute an amount of

$$\sum_{i=1}^k (-1)^{p_i} a_{d_i} 2^{i-1} \mod p,$$

where the signs are determined by the components of p. Because of the definition of k, we only have  $(-1)^{p_i} = -1$  for those positions where  $a_{d_i} = 0$ , for  $1 \leq i < k$ . But these are precisely the positions where b has ones. Hence, we find

$$val(b) - val(a) = -\sum_{i=1}^{k} 2^{i-1} + (-1)^{p_k} (b_k - a_k) 2^{k-1}.$$
 (15)

If  $a_k = p_k = 1$ , then  $b_k = 0$ , and if  $a_k = p_k = 0$ , then  $b_k = 1$ , so the second term in the rhs always equals  $2^{k-1}$ . We conclude that  $val(b) - val(a) = -(2^{k-1}-1) + 2^{k-1} = 1$ . The only exception occurs when  $a_{d_j} = p_j$  for all  $j, \leq j \leq (p-1)/2$ . In that case k is not defined. So, we proved parts (i) and (ii) under the assumption that 2 generates GF(p), which is true if and only if  $p = \pm 3 \mod 8$ , or equivalently, when  $\chi(2) = -1$ . Similar results can be obtained in the case  $p = \pm 1 \mod 8$ .

We may conclude from Theorems 3 and 4, applying eq. (6), that for all p the following result holds.

**Corollary** For all  $i \neq k$  one has  $|A_i| = N(p)$ , whereas  $|A_i| = N(p) + 1$  for  $p = \pm 1 \mod 8$ , and  $|A_i| = N(p) - 1$  for  $p = \pm 3 \mod 8$ .

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**Example** In our example p = 11, we now take m = 2. For this *m*-value,  $I_1 = \{1, 2\}$  and  $I_2 = \{3, 4, 5\}$ . The 5-cycle (13) equals  $d = (1 \ 2 \ 4 \ 3 \ 5)$ , and hence p = (0, 0, 0, 1, 0).

For  $a = (1, 1, 1, 0, 1) \in A_0$ , we find k = 3 and t = (1, 1, 1, 0, 0)Q = (1, 1, 0, 1, 1). So, b = a + t = (0, 0, 1, 1, 1), which indeed is a vector in  $A_1$ . If we take  $a = (1, 1, 0, 1, 1) \in A_2$ , then k is not defined, illustrating Theorem 4(i), since  $a_p = (1, 1, 1, 0, 1)Q = (1, 1, 0, 1, 1)$ . Taking for a the vectors ((0, 1, 0, 0, 0) and (1, 0, 1, 1, 1), both from  $A_2$ , yields (1, 1, 0, 0, 0) and (0, 1, 1, 1, 1), respectively. The third vector  $(0, 0, 1, 0, 0) \in A_3$  is the complement  $a_p$ , thus confirming Theorem 4(ii).

As an illustration of Theorem 4(iii), we consider the simple case of p = 7, where k = 3. A generator of  $GF(7)^*$  is -2. The corresponding matrix P, as defined by (8), stands for the cycle (1 2 3). Now, if we continue our construction with 2 (though 2 is not a generator), we have  $I_1 = \{1\}$  and  $I_2 = \{2,3\}$ , and therefore p = (0,0,1). Applying this vector, yields the following translations:

 $a = (0, 1, 0) \in A_2 \to (1, 1, 0) \in A_3, \ a = (0, 0, 1) \in A_3 \to (1, 0, 1) \in A_4$ 

In both translations k is equal to 1, while k is not defined for the vector  $p^c = (1, 1, 0)$ .

### 4 Remarks

Research on this topic is still in progress. Our primary motive was to develop a new approach, i.e in the context of algebraic coding theory, to the old and famous problem of determining the sign of the Gauss sum G(2) (cf. [1] for a probably exhausting list of papers on this issue). It turns out that this problem is equivalent to determining the sign of  $n_e - n_0$  (see Section 1) in the codes  $A_i$ . It was this background of which forced us to require the size of the parts in a partition not to exceed (p-1)/2. Actually, this condition is not too restrictive, since partitions of a containing one part of size (p-1)/2, can be dealt with by considering the partitions of a - (p-1)/2 as defined in this paper. Theorems 1 and 2 have their origin in [2, Lemma 4.2.4.4].

## References

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