

# The nonexistence of the (23,13,2,2) superimposed codes<sup>1</sup>

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**Abstract.** The nonexistence of (23,13,2,2) superimposed codes is proved.

## 1 Introduction

**Definition 1** A binary  $N \times T$  matrix  $C = (c_{ij})$  is called an  $(N, T, w, r)$  superimposed code (SIC) if for any pair of subsets  $W, R \subset \{1, 2, \dots, T\}$  such that  $|W| = w$ ,  $|R| = r$  and  $W \cap R = \emptyset$ , there exists a coordinate  $i \in \{1, 2, \dots, N\}$  such that  $c_{ij} = 1$  for all  $j \in W$  and  $c_{ij} = 0$  for all  $j \in R$ .

Let  $N(T, w, r)$  be the minimum length  $N$  for which an  $(N, T, w, r)$  SIC exists for fixed values of  $T$ ,  $w$  and  $r$ . The problem of determining the exact values of  $N(T, w, r)$  is completely solved only for  $w = r = 1$  [6].

The exact values of  $N(T, 2, 2)$  are known only for  $T \leq 12$  [1], [3], [2]:

$T$	4	5	6	7	8	9	10	11	12
$N(T, 2, 2)$	6	10	14	14	14	18	20	22	22

A (16,26,2,2) superimposed code is constructed in [3], hence  $22 \leq N(T, 2, 2) \leq 26$  for  $T = 13, 14, 15, 16$ .

The main result of this article is that there is no (23,13,2,2) superimposed code. Consequently  $24 \leq N(13, 2, 2) \leq 26$  for  $T = 13, 14, 15, 16$ .

## 2 Preliminaries

For a binary matrix or vector  $C$  denote by  $wt(C)$  the number of 1's in  $C$ , and by  $wt(\overline{C})$  the number of 0's in  $C$ .

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<sup>1</sup>Partially supported by the Technical University of Gabrovo under Grant C-801/2008.

For a binary matrix  $C$  denote by  $d(x, y)$  the Hamming distance between two columns  $x$  and  $y$  of  $C$ . Let  $d_2 = \min\{d(x, y) \mid x, y \in C, x \neq y\}$  and  $d(C) = \sum_{x, y \in C, x \neq y} d(x, y)$ .

**Lemma 2** (*Plotkin bound*) [5]  $\binom{T}{2} d_2 \leq N \left\lfloor \frac{T}{2} \right\rfloor \left\lfloor \frac{T+1}{2} \right\rfloor$ .

**Definition 3** Let  $x$  be a column in the superimposed code  $C$ . The residual code  $Res(C, x = a)$  is the code obtained in the following way:

- 1) take the  $i$ -th row ( $i = 1, 2, \dots, N$ ) iff  $c_{i,x} = a$ ;
- 2) delete the column  $x$  in the selected rows.

We will use the shorter notation  $Res(C, x = a, y = b)$  instead of  $Res(Res(C, x = a), y = b)$ .

**Lemma 4** Suppose  $C$  is an  $(N, T, w, r)$  superimposed code with  $w > 1, r > 1$ . Then

- (a)  $N(T-1, w-1, r) \leq wt(x) \leq N - N(T-1, w, r-1)$  for any column  $x$ ;
- (b)  $d_2 \geq 2N(T-2, w-1, r-1)$ .

*Proof.* (a) The residual code  $Res(C, x = 1)$  is a  $(wt(x), T-1, w-1, r)$  SIC, while  $Res(C, x = 0)$  is a  $(wt(\bar{x}), T-1, w, r-1)$  SIC. Hence  $wt(x) \geq N(T-1, w-1, r)$  and  $wt(\bar{x}) \geq N(T-1, w, r-1)$ .

(b) Let  $x$  and  $y$  be an arbitrary pair of different columns of  $C$ . The residual codes  $Res(C, x = 1, y = 0)$  and  $Res(C, x = 0, y = 1)$  are  $(N', T-2, w-1, r-1)$  and  $(N'', T-2, w-1, r-1)$  SIC, respectively. Hence  $d(x, y) = N' + N'' \geq 2N(T-2, w-1, r-1)$ .  $\square$

**Lemma 5** Let  $x$  be a column of  $C$  and  $A = Res(C, x = 1)$ . Then

$$d(A) + wt(\bar{A}) \geq \binom{T}{2} d_2 - (N - wt(x)) \left\lfloor \frac{T}{2} \right\rfloor \left\lfloor \frac{T+1}{2} \right\rfloor$$

*Proof.* Denote by  $C_1$  the submatrix of  $C$ , containing all rows with value 1 in the column  $x$ , and by  $C_0$  the remaining part of  $C$ . Then

$$\binom{T}{2} d_2 \leq d(C) = d(C_1) + d(C_0).$$

But  $d(C_1) = d(A) + wt(\bar{A})$  and  $d(C_0) \leq (N - wt(x)) \left\lfloor \frac{T}{2} \right\rfloor \left\lfloor \frac{T+1}{2} \right\rfloor$ .

The result follows.  $\square$

**Definition 6** Two  $(N, T, w, r)$  superimposed codes are equivalent if one of them can be obtained from the other by a permutation of the rows and a permutation of the columns. In the case  $w = r$  an inversion of the all code entries is also allowed.

### 3 Main result

**Lemma 7** *If  $C$  is a  $(23, 13, 2, 2)$  superimposed code then*

(a)  $9 \leq wt(x) \leq 14$  for any column  $x$  of  $C$ ;

(b)  $d_2 = 12$ .

*Proof.* (a) Follows from Lemma 4 and the known value  $N(12, 1, 2) = 9$  [3];

(b) Follows from Lemma 2, Lemma 4 and the known value  $N(11, 1, 1) = 6$  [6].  $\square$

**Theorem 8** *There is no  $(23, 13, 2, 2)$  superimposed code.*

*Proof.* Suppose  $C$  is a  $(23, 13, 2, 2)$  SIC. Up to equivalence we may assume that the code has the following form:

$$\left( \begin{array}{c|c} 1 & \\ \vdots & A \\ 1 & \\ \hline 0 & \\ \vdots & B \\ 0 & \end{array} \right)$$

where the matrix  $A$  is a  $(N, 12, 1, 2)$  SIC where  $N \in \{9, 10, 11\}$ , and the matrix  $B$  has to be chosen in such a way that the whole matrix to be a  $(23, 13, 2, 2)$  SIC. We may assume that the rows of  $B$  are sorted lexicographically.

Applying the method described in [4] we constructed all inequivalent  $(9, 12, 1, 2)$ ,  $(10, 12, 1, 2)$  and  $(11, 12, 1, 2)$  superimposed codes. Then we checked the condition of Lemma 5, which turned out to reduce the amount of computations approximately 6 times.

SIC parameters	$(9, 12, 1, 2)$	$(10, 12, 1, 2)$	$(11, 12, 1, 2)$
number of inequivalent SIC	1	99	243709
number of inequivalent SIC, which satisfy Lemma 5	1	54	44509

Using an exhaustive computer search we tried to construct the matrix  $B$  column by column taking into account the restrictions of Lemma 7 and the sorted rows property. It turned out, however, that superimposed codes with parameters  $(23, 13, 2, 2)$  do not exist.  $\square$

## References

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