# On the (2,1)-extendability of ternary linear codes

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#### Overview

The (2,1)-extendability of  $[n, k, d]_3$  codes with gcd(3, d) = 1 is investigated geometrically for  $3 \le k \le 5$ .

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#### **0.** Introduction

C:  $[n, k, d]_3$  code,  $k \ge 3$ , gcd(d, 3) = 1. The diversity  $(\Phi_0, \Phi_1)$  of C is given as

$$\Phi_0 = \frac{1}{2} \sum_{\substack{3 \mid i, i \neq 0}} A_i, \quad \Phi_1 = \frac{1}{2} \sum_{\substack{i \neq 0, d \pmod{3}}} A_i.$$

$$\begin{split} \mathcal{D}_k &= \text{the set of all possible diversities of} \\ & [n,k,d]_3 \text{ codes.} \\ \mathcal{D}_k^* &= \{(\theta_{k-2},0),(\theta_{k-3},2\cdot 3^{k-2}),\\ & (\theta_{k-2},2\cdot 3^{k-2}),(\theta_{k-2}+3^{k-2},3^{k-2})\},\\ & (\theta_{k-2},2\cdot 3^{k-2}),(\theta_{k-2}+3^{k-2},3^{k-2})\},\\ \text{where } \theta_j &= (3^{j+1}-1)/2.\\ \mathcal{D}_k^+ &= \mathcal{D}_k \setminus \mathcal{D}_k^*. \end{split}$$

**Theorem 1** (Maruta, 2005). C is extendable if  $(\Phi_0, \Phi_1) \in \mathcal{D}_k^*$ .

#### Probrem.

Is  $\mathcal{C}$  (2,1)-extendable when  $\mathcal{C}$  is not extendable?

We consider this problem for  $3 \le k \le 5$ . (See Proceedings for k = 6)

# Diversities of non-extendable ternary codeskdiversity3(4,3)4(13,9), (10,15), (16,12)5(40,27), (31,45), (40,36), (40,45), (49,36)6(121,81), (94,135), (121,108), (112,126), (130,117), (121,135), (148,108)

How about (2, 1)-extendability of such codes?

#### 1. Geometric approach

 $\mathcal{C}$ :  $[n, k, d]_3$  code,  $k \geq 3$  $G = [g_1, \cdots, g_k]^{\mathsf{T}}$ : a generator matrix of  $\mathcal{C}$  $\Sigma := PG(k-1,3)$ : the projective space of dimension k-1 over  $\mathbb{F}_3$ For  $P = P(p_1, \ldots, p_k) \in \Sigma$  we define the weight of P with respect to C, denoted by  $w_{\mathcal{C}}(P)$ , as  $w_{\mathcal{C}}(P) = wt(\sum_{i=1}^{k} p_i g_i).$ 

Let

$$F_d = \{P \in \Sigma \mid w_{\mathcal{C}}(P) = d\},$$
  

$$F_0 = \{P \in \Sigma \mid w_{\mathcal{C}}(P) \equiv 0 \pmod{3}\},$$
  

$$F_1 = \{P \in \Sigma \mid w_{\mathcal{C}}(P) \not\equiv 0, d \pmod{3}\},$$
  

$$F_2 = \{P \in \Sigma \mid w_{\mathcal{C}}(P) \equiv d \pmod{3}\} (\supset F_d),$$
  

$$F = F_0 \cup F_1, F_e = F_2 \setminus F_d.$$

Then  $(\Phi_0, \Phi_1) = (|F_0|, |F_1|).$ 

2. (2,1)-extendability of ternary linear codes

Lemma 2. C is (2, 1)-extendable  $\Leftrightarrow \exists (k-3)$ -flat  $\subset F \cup F_e$  $\Leftarrow \exists (k-3)$ -flat  $\subset F = F_0 \cup F_1$ .

Every non-extendable  $[n, 3, d]_3$  code C with gcd(3, d) = 1 has diversity (4, 3). C is (2, 1)-extendable since  $F \neq \emptyset$ .  $\begin{aligned} \Pi_t : \text{ a } t\text{-flat in } \Sigma. \\ \text{The } \textit{diversity of } \Pi_t \text{ is } (\varphi_0^{(t)}, \varphi_1^{(t)}) \\ \text{ where } \varphi_s^{(t)} &= |\Pi_t \cap F_s|, \ s = 0, 1. \\ \Pi_t \text{ is called a } (\varphi_0^{(t)}, \varphi_1^{(t)})_t \text{ flat.} \\ (i, j)_1 \text{ flats, } (i, j)_2 \text{ flats, } (i, j)_3 \text{ flats are called} \\ (i, j)\text{-lines, } (i, j)\text{-planes, } (i, j)\text{-solids, respectively.} \end{aligned}$ 

 $c_{i,j}^{(t)}$ : the number of  $(i,j)_{t-1}$  flats in  $\Pi_t$ . The list of  $c_{i,j}^{(t)}$ 's is called its *spectrum*.

Table 1.									
$\varphi_0^{(2)}$	$\varphi_1^{(2)}$	$c_{1,0}^{(2)}$	$c_{0,2}^{(2)}$	$c_{2,1}^{(2)}$	$c_{1,3}^{(2)}$	$c_{4,0}^{(2)}$			
4	0	12	0	0	0	1			
1	6	2	9	0	2	0			
4	3	4	3	6	0	0			
4	6	0	3	6	4	0			
7	3	1	0	9	1	2			
4	9	0	0	0	12	1			
13	0	0	0	0	0	13			

 $\exists (i, j)$ -line  $\Rightarrow \exists (3i + 1, 3j)$ -plane. (4, 3) and (4, 6) are new.

Table 2.								
$\varphi_0^{(3)}$	$\varphi_1^{(3)}$	$(3) c_{4,0}^{(3)}$	$c_{1,6}^{(3)}$	$(3) c_{4,3}$	$^{(3)}_{c_{4,6}}$	$^{(3)}_{c_{7,3}}$	$(3) c_{4,9}$	$c_{13,0}^{(3)}$
13	0	39	0	0	0	0	0	1
4	18	2	36	0	0	0	2	0
13	9	4	3	27	0	6	0	0
10	15	0	10	15	15	0	0	0
16	12	0	0	12	12	16	0	0
13	18	0	3	0	27	6	4	0
22	9	1	0	0	0	36	1	2
13	27	0	0	0	0	0	39	1
40	0	0	0	0	0	0	0	40







**Theorem 3.** Let C be an  $[n, k, d]_3$  code with k = 3 or 4, gcd(3, d) = 1. Then C is (2, 1)-extendable.

**Theorem 4.** Let C be an  $[n, 5, d]_3$  code with diversity  $(\Phi_0, \Phi_1) \neq (40, 36)$ , gcd(3, d) = 1. Then C is (2, 1)-extendable.

Table 3.

$\varphi_0^{(4)}$	$\varphi_1^{(4)}$	$c_{13,0}^{(4)}$	$c_{4,18}^{(4)}$	$c_{13,9}^{(4)}$	$c_{10,15}^{(4)}$	$c_{16,12}^{(4)}$	$c_{13,18}^{(4)}$	$c_{22,9}^{(4)}$	$c_{13,27}^{(4)}$	$c_{40,0}^{(4)}$
40	0	120	0	0	0	0	0	0	0	1
13	54	2	117	0	0	0	0	0	2	0
40	27	4	3	108	0	0	0	6	0	0
31	45	0	10	15	81	0	15	0	0	0
40	36	0	0	40	36	45	0	0	0	0
40	45	0	0	0	36	45	40	0	0	0
49	36	0	0	12	0	81	12	16	0	0
40	54	0	3	0	0	0	108	6	4	0
67	27	1	0	0	0	0	0	117	1	2
40	81	0	0	0	0	0	0	0	120	1
121	0	0	0	0	0	0	0	0	0	121

**Theorem 5.** Let C be an  $[n, 5, d]_3$  code with diversity (40, 36), gcd(3, d) = 1. Then C is (2, 1)-extendable iff either (a)  $\exists l:(1, 0)$ -line,  $l \subset F_0 \cup F_e$ , or (b)  $\exists R_1, R_2, R_3 \in F_e$  s.t.  $\langle R_i, R_j \rangle$ :(0, 2)-line.



**Example.** Let  $\mathcal{C}$  be a  $[15, 5, 8]_3$  code with a generator matrix



whose weight distribution is  $0^{1}8^{60}9^{40}10^{62}11^{20}12^{40}13^{10}14^{10}$  (diversity (40,36)).



 $w_{\mathcal{C}}(R_1) = 11, \ w_{\mathcal{C}}(R_2) = 14, \ w_{\mathcal{C}}(R_3) = 11.$  $w_{\mathcal{C}}(R_i) \equiv 2 \equiv d = 8 \pmod{3}.$  $w_{\mathcal{C}}(Q_1) = 10, \ w_{\mathcal{C}}(Q_2) = 10 \cdots w_{\mathcal{C}}(Q_j) \equiv 1 \pmod{3}.$  Since  $V(x_0 + 2x_1 + 2x_2 + x_3) \cap V(x_0 + x_2 + 2x_4) = \langle R_1, R_2, R_3 \rangle$ , by adding the columns  $(1, 2, 2, 1, 0)^T$  and  $(1, 0, 1, 0, 2)^T$  to G, we get a (2, 1)-extension of C whose weight distribution is

 $0^{1}9^{38}10^{56}11^{46}12^{34}13^{30}14^{26}15^{8}16^{4}$ .

#### Remark.

We have no example of an  $[n, 5, d]_3$  code with diversity (40,36), gcd(3, d) = 1, which is not (2,1)-extendable. Theorem 6. Let C be an  $[n, 5, d]_3$  code with diversity (40, 36), gcd(3, d) = 1. Then C is (2, 1)-extendable if  $A_d \leq 50$ .

To prove this, we need: (1) In  $(40,36)_4$  flat, any (10,15)-solid has the unique focal point from  $F_1$  and any (16,12)solid has the unique focal point from  $F_2$ . (2) Every 20-cap in PG(4,3) is either a  $\Gamma$ -cap or  $\Delta$ -cap.

#### 3. Focal points and focal hyperplanes

For i = 1, 2, a point  $P \in F_i$  is called a *focal point* of a hyperplane H (or P is *focal to* H) if the following three conditions hold:

Such a hyperplane H is called a *focal hyper*plane of P (or H is *focal to* P).

#### Theorem 7

Let  $\Pi_t$  be a *t*-flat with new diversity. Then, for i = 1, 2, any point of  $\Pi_t \cap F_i$  has a focal (a, b)-hyperplane in  $\Pi_t$  for some new (a, b). **Note.** 

Let (2,1) and (0,2) be new in the types of lines for convenience.









**Lemma 8.** Let  $\delta$  be a (4,3)-plane. Then, every point of  $\delta \cap F_1$  and of  $\delta \cap F_2$  has a focal (0,2)-line and a focal (2,1)-line, respectively, and vice versa.





**Lemma 9.** Let  $\delta$  be a (4,6)-plane. Then, every point of  $\delta \cap F_1$  and of  $\delta \cap F_2$  has a focal (2,1)-line and a focal (0,2)-line, respectively, and vice versa. Lemma 10. Let  $\delta$  be a (16, 12)-solid. Then, every point of  $\delta \cap F_1$  and of  $\delta \cap F_2$  has a focal (4,3)-plane and a focal (4,6)-plane, respectively, and vice versa.



**Lemma 11.** Let  $\Pi$  be a  $(40, 36)_4$  flat. Then, every point of  $\Pi \cap F_1$  and of  $\Pi \cap F_2$  has a focal (10, 15)-solid and a focal (16, 12)-solid, respectively, and vice versa.

#### 4. Proof of Theorem 6.

Assume C is not (2, 1)-extendable. Then no three points of  $F_e$  are collinear by Theorem 5.

 $F_e$  forms a cap and we have  $|F_e| \leq 20$  since the largest size of a cap in PG(4,3) is 20.

Every 20-cap in PG(4,3) is either a  $\Gamma$ -cap or a  $\Delta$ -cap (Pellegrino, 1974).





 $\Delta\text{-}cap = \{\bullet\} \cup \{\bullet\} \cup \{\bullet\} \cup \{\bullet\} \cup \{\bullet\}$ 

We only consider the case when  $F_e$  is a 20cap of type  $\Gamma$ . (See Proceedings for the case when  $F_e$  is a 20-cap of type  $\Delta$ .)



## Case 1: $V \in F_0 \Rightarrow VP_i = (1, 0)$ -line

![](_page_39_Figure_1.jpeg)

Actually, there are exactly six (1, 0)-lines through V in the  $(40, 36)_4$  flat  $\Sigma$ , a contradiction.

![](_page_40_Figure_0.jpeg)

H:solid

#### Case 2: $V \in F_1 \Rightarrow VP_i = (0, 2)$ -line

![](_page_41_Figure_1.jpeg)

H:solid

![](_page_42_Figure_0.jpeg)

## H': the focal (10, 15)-solid of V

E': the projection of E from V onto H' if  $H \neq H'$ .

### V is point of $F_1 \Rightarrow VP_i = (0, 2)$ -line

![](_page_43_Figure_1.jpeg)

In *H'*, at most  $c_{1,6}^{(3)} + c_{4,6}^{(3)} = 25$  planes meet *E'* in four points, a contradiction.

#### Case 3: $V \in F_2 \Rightarrow VP_i = (1, 0)$ -line

![](_page_44_Figure_1.jpeg)

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# Thank you for your attention!